On the $q$-criticality of graphs with respect to secure graph domination

Authors’ identities suppressed: Blind refereeing copy

Abstract

A subset $X$ of the vertex set of a graph $G$ is a secure dominating set of $G$ if each vertex of $G$ which is not in $X$ is adjacent to some vertex in $X$ and if, for each vertex $u$ not in $X$, there is a neighbouring vertex $v$ of $u$ in $X$ such that the swap set $(X - \{v\}) \cup \{u\}$ is again a dominating set of $G$. The secure domination number of $G$ is the cardinality of a smallest secure dominating set of $G$.

The notion of secure graph domination finds applications in the generic setting where the vertex set of $G$ represents a distributed spatial domain of locations and the edges of $G$ represent links between these locations. Patrolling guards, each stationed at one of these locations, may move along the links in order to protect the graph. A minimum secure dominating set of $G$ then represents a smallest collection of locations at which guards may be stationed so that the entire location complex modelled by $G$ is protected in the sense that if a security concern arises at location $u$, there will either be a guard stationed at that location who can deal with the problem, or else a guard dealing with the problem from an adjacent location $v$ will still leave the location complex protected after moving from location $v$ to location $u$ in order to deal with the problem.

A graph $G$ is $q$-critical if the smallest arbitrary subset of edges whose removal from $G$ necessarily increases the secure domination number, has cardinality $q$. The notion of $q$-criticality is important in applications such as the one mentioned above, because it provides threshold information as to the number of edge failures (perhaps caused by an adversary) that will necessitate the hiring of additional guards to secure the location complex. Denote by $\Omega_n$ the largest value of $q$ for which $q$-critical graphs of order $n$ exist. It has previously been established that $\Omega_2 = 1$, $\Omega_3 = 2$, $\Omega_4 = 4$, $\Omega_5 = 6$ and $\Omega_6 = 9$. In this paper we present a repository of all $q$-critical graphs of orders $2, 3, 4, 5$ and $6$ for all admissible values of $q$ and we also establish the previously unknown values $\Omega_7 = 12$ and $\Omega_8 = 17$. These values support an existing conjecture that $\Omega_n = (\frac{n}{2}) - 2n + 5$ for all $n \geq 7$.

Key words: Secure domination, graph protection, edge criticality.

1 Introduction

A dominating set of a graph $G$ is a subset $X$ of the vertex set of $G$ with the property that each vertex of $G$ not in $X$ is adjacent to at least one vertex in $X$. A secure dominating set of $G$ is a subset $X_s$ of the vertex set of $G$ with the property that $X_s$ forms a dominating set of $G$ and, additionally, for each vertex $u$ not in $X_s$, that there exists a vertex $v \in X_s$.
for which the swap set \((X_s − \{v\}) ∪ \{u\}\) is again a dominating set of \(G\). The secure domination number of \(G\), denoted by \(\gamma_s(G)\), is the minimum value of \(|X_s|\), taken over all secure dominating sets \(X_s\) of \(G\) (i.e. the cardinality of a smallest secure dominating set of \(G\)). A number of general bounds have been established for the parameter \(\gamma_s(G)\) in [7], and exact values of \(\gamma_s(G)\) have also been established for various graph classes, such as paths, cycles, complete multipartite graphs and products of paths and cycles. Various properties of secure dominating sets of graphs have also been studied in [1, 2, 4, 5, 6].

Consider, as an example, the graph \(G_1\) in Figure 1, \(\gamma_s(G_1) = 2\). A minimum secure dominating set for \(G_1\) is \({v_1, v_2}\); vertex \(v_4\) is defended by \(v_2\) while \(v_3\) and \(v_5\) are both defended by \(v_1\).

\[\text{Figure 1: A minimum secure dominating set } \{v_1, v_2\} \text{ for a graph } G_1 \text{ of order } 5.\]

The notion of secure graph domination finds applications in the generic setting where the vertex set of \(G\) represents a distributed spatial domain of locations and the edges of \(G\) represent links between these locations. Patrolling guards, each stationed at one of these locations, may move along the links in order to protect the graph. A minimum secure dominating set of \(G\) then represents a smallest collection of locations at which guards may be stationed so that the entire location complex modelled by \(G\) is protected in the sense that if a security concern arises at location \(u\), there will either be a guard stationed at that location who can deal with the problem, or else a guard dealing with the problem from an adjacent location \(v\) will still leave the location complex protected after moving from location \(v\) to location \(u\) in order to deal with the problem.

In this setting, the notion of edge removal is important, because one might seek the cost (in terms of the additional number of guards required to protect a location complex modelled by \(G\)) if a number of edges in \(G\) fail (i.e. a number of links are eliminated form the location complex, thereby disqualifying guards from moving along such disabled links).

A graph \(G\) is \(q\)-critical if the smallest arbitrary subset of edges whose removal from \(G\) necessarily increases the secure domination number, has cardinality \(q\). Being able to determine the value of \(q\) for which a given graph is \(q\)-critical is important from an application point of view, because this value may be seen as a robustness threshold in the sense that the failure of any \(q − 1\) edges in \(G\) results in a graph that can still be dominated securely by \(\gamma_s(G)\) guards, but this is not true for the failure of any set of \(q\) edges in \(G\).

In this paper, we provide empirical evidence in support of a conjecture by Burger et al. [3] that the largest value of \(q\) for which there exists a graph of order \(n\) that is \(q\)-critical, is \(\binom{n}{2} − 2n + 5\) for all \(n ≥ 7\). We also provide a repository of all \(q\)-critical graphs of order \(n\) and size \(m\) for all \(q ∈ \{0, 1, \ldots, m\}\) and all \(m ∈ \{1, 2, \ldots, \binom{n}{2}\}\), where \(n ∈ \{2, \ldots, 6\}\).
The concept of $q$-criticality

We denote the set of all non-isomorphic graphs obtained by removing $q \in \{0, 1, \ldots, m\}$ edges from a given graph $G$ of size $m$ by $G - qe$. Furthermore, let $\gamma_s(G - qe)$ denote the set of values of $\gamma_s(H)$ as $H \in G - qe$ varies (for a fixed value of $q$). We distinguish between the graph obtained by removing a specific edge $e$ from a given graph $G$, by writing $G - e$, and the class of graphs obtained by removing any single edge from $G$, by writing $G - 1e$.

The cost function

$$c_q(G) = \min \gamma_s(G - qe) - \gamma_s(G)$$  \hspace{1cm} (1)

is nonnegative and bounded from above by $q$ for all $q \in \{0, 1, \ldots, m\}$ [3]. This cost function measures the smallest possible increase in the secure domination number of a member of $G - qe$, relative to the secure domination number of a graph $G$ of size $m$, when a set of $q \in \{0, 1, \ldots, m\}$ edges are removed from $G$.

![Figure 2: The edge-removal metagraph of the complete graph $K_4$ of order 4. The set $K_4 - qe$ is shown on level $q$ of the graph for all $q = 0, \ldots, 6$. Minimum secure dominating sets of the resulting graphs are denoted by solid vertices in each case. An arrow of the form $G \rightarrow H$ from level $q$ to level $q + 1$ means that $G$ is a certificate in $K_n - qe$ showing that $H \in K_n - (q + 1)e.$]
A graph $G$ is $q$-critical if $c_{q-1}(G) = 0$, but $c_q(G) > 0$ (that is, if the smallest arbitrary subset of edges whose removal from $G$ necessarily increases the secure domination number, has cardinality $q$). The notion of $q$-criticality partitions the set of all non-isomorphic, nonempty graphs of order $n$ in the sense that any such graph is $q$-critical for exactly one value of $q \in \{0, 1, 2, \ldots, \binom{n}{2}\}$, as demonstrated in the so-called edge-removal graph of the complete graph $K_n$ of order $n$ in Figure 2. The edge-removal metagraph of a graph $G$ of size $m$ is a graph whose nodes represent the non-isomorphic members of $G - e$ for all $q = 0, 1, \ldots, m$. The nodes on level $q$ correspond to the members of $G - e$. A node $H$ on level $q - 1$ of this metagraph is joined to a node $H'$ on level $q$ if $H'$ can be obtained by removing one edge from $H$, for any $q \in \{1, 2, \ldots, m\}$. The only node on level 0 of the edge-removal metagraph of some graph $G$ corresponds to $G$ itself, while the only node on level $m$ corresponds to the empty graph of the same order as $G$. The edge-removal metagraph of $G$, and the edge-removal metagraph of the complete graph $K_n$ is of particular interest, because it contains nodes corresponding to all the non-isomorphic graphs of order $n$.

Let $Q^q_n$ be the class of $q$-critical graphs of order $n \geq 2$ and some $q \in \{0, 1, 2, \ldots, \binom{n}{2}\}$. Grobler and Mynhardt characterised the graph class $Q^1_n$ for all $n \in \mathbb{N}$ in 2009 [8, Theorem 2] and used their characterisation to derive a four-step construction process for computing all the members of the class $Q^q_n$. Because of space constraints we do not give a full description of this (nontrivial) construction process here, but rather refer the reader to [8, Section 3.1] for the details. The following characterisation may be used to compute the class $Q^q_n$ inductively from the class $Q^{q-1}_n$ for any integer $n \geq 2$ and all permissible values of $q \geq 2$, using the above-mentioned 4-step construction process by Mynhardt and Grobler [8] for the class $Q^1_n$ as base case.

**Proposition 1 ([3])** A graph $G$ of size at least $q > 1$ is $q$-critical if and only if 

(a) at least one graph $H^* \in G - e$ for which $\gamma_s(H^*) = \gamma_s(G)$ is $(q - 1)$-critical, and 

(b) each graph $H \in G - e$ for which $\gamma_s(H) = \gamma'_s(G)$ is $q'$-critical for some $q' \leq q - 1$. ■

The inductive process referred to above is formalised in Algorithm 1. The algorithm commences by considering a graph $H \in Q^{q-1}_n$ and proceeding to add a single edge $e \notin E(H)$ to $H$ in Step 3, upon which the result of Proposition 1 is used to test whether or not $H + e \in Q^q_n$. This process is repeated for each edge $e \notin E(H)$ and for each graph $H \in Q^{q-1}_n$.

**Algorithm 1: Computing the class of $Q^q_n$ of $q$-critical graphs of order $n$**

```
Input : The graph classes $Q^1_n, \ldots, Q^{q-1}_n$.
Output : The class $Q^q_n$ of $q$-critical graphs of order $n$.
1 for each $H \in Q^{q-1}_n$ do
2   for each $e \notin E(H)$ do
3     if $q$-Critical$(H + e, q)$ then $Q^q_n \leftarrow Q^q_n \cup \{H + e\}$
```

In Step 3 of Algorithm 1, another algorithm, Algorithm 2, is called to test whether $G = H + e \in Q^q_n$. In Algorithm 2, each member of $G - e$ is examined. If a member $E \in G - e$
On the q-criticality of graphs with respect to secure graph domination

is found for which \( \gamma_s(E) = \gamma_s(G) \), then \( G \notin Q^0_n \) by Proposition 1. Similarly, if a member \( F \in Q^p_n \) is found for some \( p \geq q \), then \( G \notin Q^q_n \) by Proposition 1. If, however, no such graph \( F \) is found, then \( G \in Q^q_n \) by Proposition 1, since \( H \in Q^{q-1}_n \).

Algorithm 2: q-Critical\((G,q)\)

\[
\text{Input : A graph } G \text{ and the value of } q. \\
\text{Output : A boolean value stating whether } G \text{ is } q\text{-critical.}
\]

1. if \( G \in Q^p_n \) for some \( p \leq q - 1 \) then
2. return [FALSE]
3. for each \( e \in E(G) \) do
4. if \( \gamma_s(G - e) = \gamma_s(G) \) and \( G - e \notin Q^p_n \) for some \( p \leq q - 1 \) then
5. return [FALSE]
6. return [TRUE]

Figure 3: The graph classes \( Q^1_1 \), \( Q^2_1 \), \( Q^3_1 \) and \( Q^4_1 \). Minimum secure dominating sets are denoted by solid vertices in each case. An arrow of the form \( H^* \to G \) in the figure denotes the relationship between the graphs \( G \) and \( H^* \) in Proposition 1.

The graph classes \( Q^1_1, \ldots, Q^4_1 \) are shown in Figure 3. The classes \( Q^3_1 \) and \( Q^4_1 \) are both empty. The 4-step construction of Grobler and Mynhardt [8] was used to compute the class \( Q^1_n \) in the first column of the figure as base case. Thereafter, Algorithm 1 was used to compute the classes \( Q^2_1, Q^3_1 \) and \( Q^4_1 \) inductively.

Note that it is, in view of Proposition 1 and Algorithms 1–2, not necessary to construct the entire edge-removal metagraph of the complete graph of order \( n \) in order to determine the graph class \( Q^q_n \) for a fixed value of \( q \); instead only the classes \( Q^1_n, \ldots, Q^q_n \) need be constructed inductively which, for values of \( q \) that are small compared to \( n \), can be achieved.
Authors' identities suppressed: Blind refereeing copy

<table>
<thead>
<tr>
<th>$n \to$</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
</tr>
</thead>
<tbody>
<tr>
<td>$\Omega_n \to$</td>
<td>1</td>
<td>2</td>
<td>4</td>
<td>7</td>
<td>14</td>
<td>26</td>
<td>52</td>
</tr>
<tr>
<td>$</td>
<td>Q_1^n</td>
<td>$</td>
<td>1 (≪ 1)</td>
<td>2 (≪ 1)</td>
<td>6 (&lt; 1)</td>
<td>18 (2)</td>
<td>50 (10)</td>
</tr>
<tr>
<td>$</td>
<td>Q_2^n</td>
<td>$</td>
<td>3 (≪ 1)</td>
<td>9 (&lt; 1)</td>
<td>32 (3)</td>
<td>111 (22)</td>
<td>428 (87)</td>
</tr>
<tr>
<td>$</td>
<td>Q_3^n</td>
<td>$</td>
<td>1 (&lt; 1)</td>
<td>8 (&lt; 1)</td>
<td>34 (4)</td>
<td>165 (48)</td>
<td>911 (326)</td>
</tr>
<tr>
<td>$</td>
<td>Q_4^n</td>
<td>$</td>
<td>1 (&lt; 1)</td>
<td>18 (6)</td>
<td>195 (102)</td>
<td>1878 (2265)</td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_5^n</td>
<td>$</td>
<td>8 (1)</td>
<td>153 (92)</td>
<td>2012 (3931)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_6^n</td>
<td>$</td>
<td>2 (1)</td>
<td>93 (89)</td>
<td>1848 (5222)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_7^n</td>
<td>$</td>
<td>1 (&lt; 1)</td>
<td>37 (51)</td>
<td>1521 (5170)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_8^n</td>
<td>$</td>
<td>10 (23)</td>
<td>1089 (4378)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_9^n</td>
<td>$</td>
<td>3 (10)</td>
<td>627 (2852)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{10}^n</td>
<td>$</td>
<td>1 (3)</td>
<td>260 (1335)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{11}^n</td>
<td>$</td>
<td>76 (442)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{12}^n</td>
<td>$</td>
<td>19 (134)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{13}^n</td>
<td>$</td>
<td>5 (44)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{14}^n</td>
<td>$</td>
<td>2 (15)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$</td>
<td>Q_{15}^n</td>
<td>$</td>
<td>1 (7)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Graph Total</td>
<td>1</td>
<td>3</td>
<td>10</td>
<td>33</td>
<td>155</td>
<td>1043</td>
<td>12345</td>
</tr>
<tr>
<td>Time Total</td>
<td>≪ 1</td>
<td>≪ 1</td>
<td>&lt; 1</td>
<td>2</td>
<td>23</td>
<td>531</td>
<td>27208</td>
</tr>
</tbody>
</table>

Table 1: Cardinalities of the graph classes $Q_1^n, \ldots, Q_{\Omega_n}^n$ for $n \in \{2, \ldots, 8\}$ as computed on a 3.4 GHz Intel(R) Core(TM) i7-3770 processor with 8 GiB RAM running in Ubuntu 12.04 and using a C++ implementation of Algorithms 1–2 in conjunction with the Boost graph library [11] for graph isomorphism testing. Computation times, indicated in brackets, are measured in seconds and represent the time required for determining the graph class $Q_q^n$ from the graph class $Q_{q-1}^n$, for all $q \in \{2, \ldots, \Omega_n\}$.

in a fraction of the time required to construct the entire edge removal metagraph of $K_n$.

3 Numerical results

Let $\Omega_n$ denote the largest value of $q$ for which there exist $q$-critical graphs of order $n$. Values of $\Omega_n$ have been established for small $n$. In particular, Burger et al. [3] showed that $\Omega_2 = 1, \Omega_3 = 2, \Omega_4 = 4, \Omega_5 = 6$ and $\Omega_6 = 9$. They also conjectured as follows.

Conjecture 1 ([3]) $\Omega_n = \binom{n}{2} - 2n + 5$ for all $n \geq 7$.

In this paper we provide further circumstantial evidence in support of Conjecture 1, by proving the conjecture correct for $n \in \{7, 8\}$. In particular, using a C++ implementation of
Algorithms 1–2, we confirm the values of $\Omega_n$ for $n \leq 6$ mentioned above, and additionally show that $\Omega_7 = 12$ and $\Omega_8 = 17$. The results thus obtained are summarized in Table 1, which contains listings of the cardinalities of the graph classes $Q^n_q$ for $n \in \{2, \ldots, 8\}$ and $q \in \{1, \ldots, \Omega_n\}$. The classes $Q^1_n, \ldots, Q^\Omega_n$ were determined by the 4-step construction process of Mynhardt and Grobler [8], referred to above.

Table 2: The graph classes $Q^1_n, \ldots, Q^\Omega_n$ for $n \in \{2, \ldots, 6\}$. Class members are presented in the well-known graph6 format, which is ideal for storing undirected graphs in a compact manner, using only printable ASCII characters. These graph representations may be converted to adjacency matrices (or other formats) using the reader showg which is available online [9].

A repository of the members of the graph classes $Q^1_n, \ldots, Q^\Omega_n$ is provided in Table 2 for $n \in \{2, 3, 4, 5, 6\}$. The graphs in this table are presented in the well-known graph6 format [9], which is ideal for storing class representatives of isomorphism classes of undirected graphs in a compact manner, using only printable ASCII characters. These graphs may be converted to adjacency matrices and other formats using the reader showg, which is available online [9]. The reader showg package is part of nauty, originally designed by McKay and Piperno [10] for graph isomorphism testing.
4 Further work

In addition to attempting a general proof or refutation of Conjecture 1, another interesting problem for future research would be to investigate the largest number of arbitrary edges whose removal from a graph necessarily does not increase the secure domination number. In this context the cost function

\[ C_p(G) = \max \gamma_s(G - pe) - \gamma_s(G), \]

is applicable instead of (1), and a graph \( G \) may be defined to be \( p \)-stable if \( C_p(G) = 0 \), but \( C_{p+1}(G) > 0 \). This problem finds application in cases where one seeks threshold information in terms of the largest set of edges whose removal from \( G \) does not increase the secure domination number of the resulting graph.

Acknowledgements

The research towards this paper was supported financially by the South African National Research Council under grant numbers 70593, 77248 and 81558.

References


